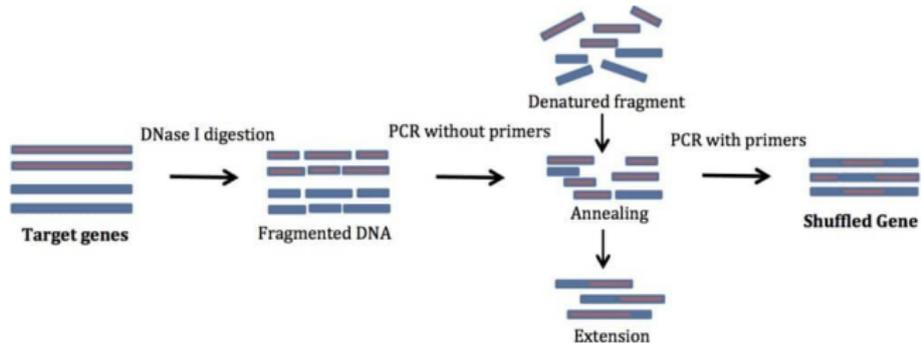
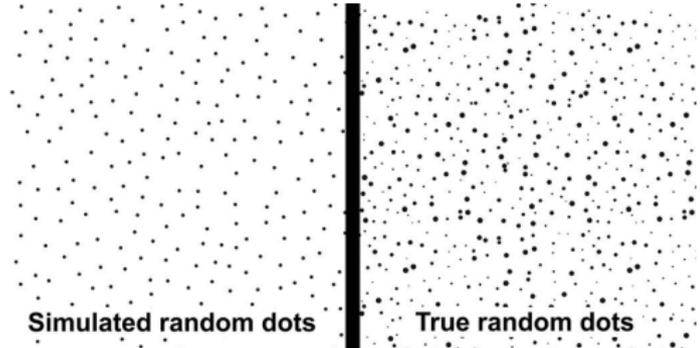


**What is...card shuffling mathematically?**

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Or: Random walks

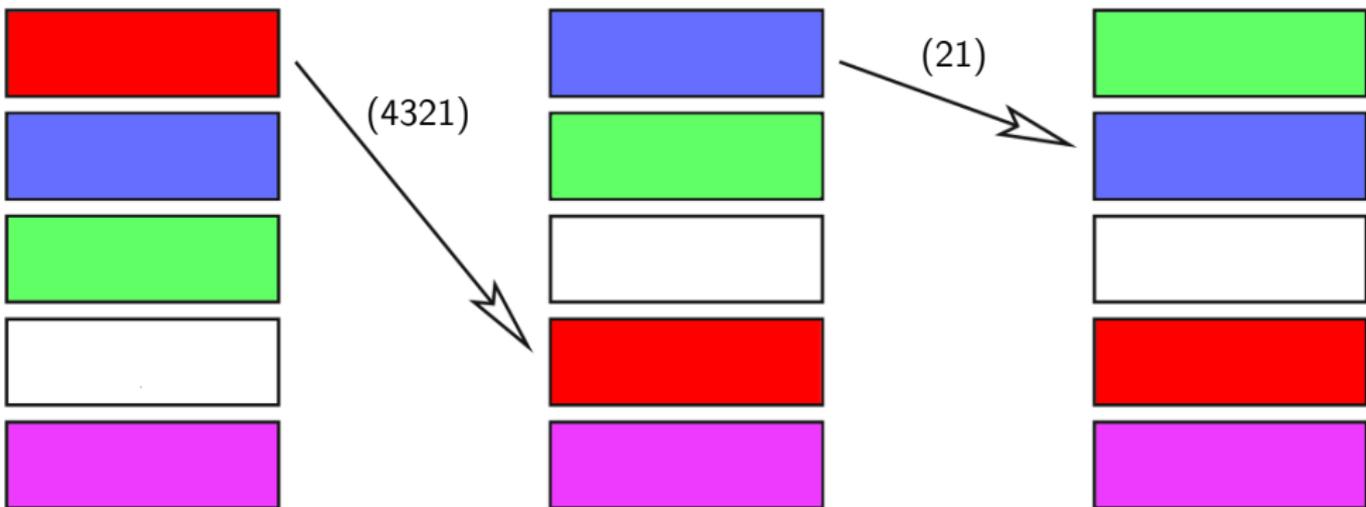
# Shuffle problems are everywhere



▶ Shuffling is everywhere: cards, diffusion, playlist, DNA, many more...

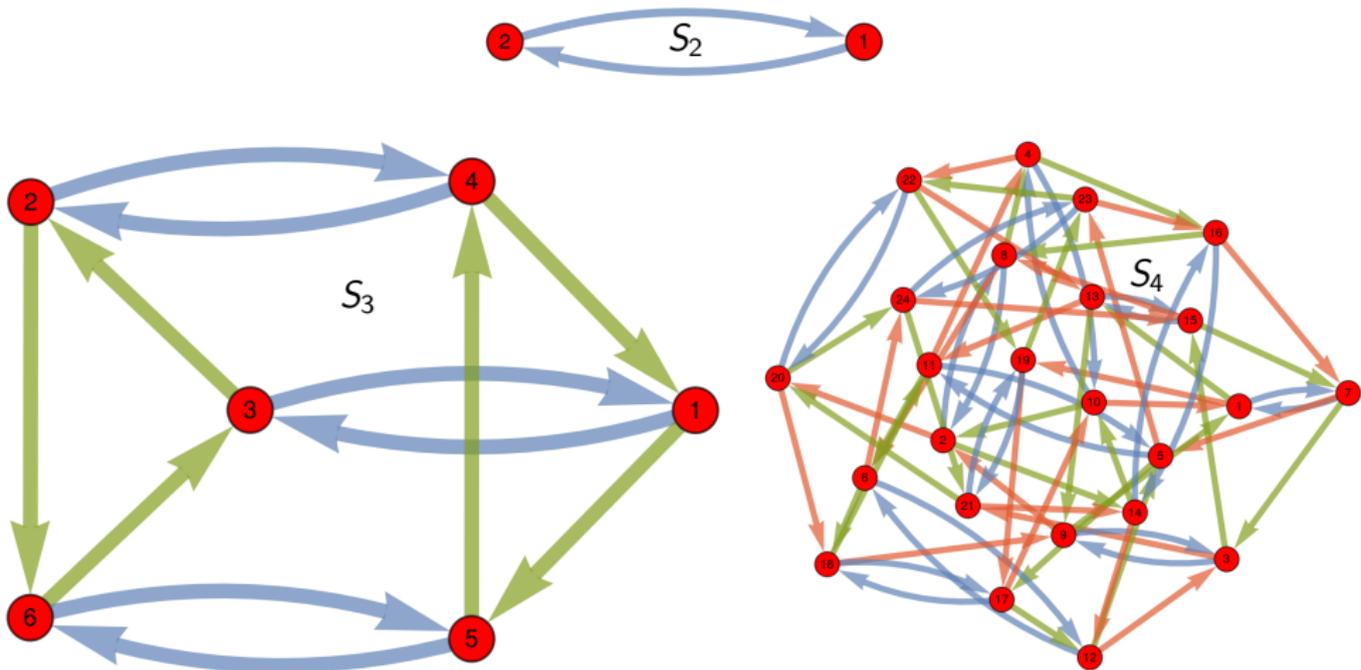
▶ **Goal** Model shuffling mathematically

## An inefficient way to shuffle



- ▶ The shuffle we want to analyze is **top-to-random** of  $n$  cards
- ▶ Every shuffle is modeled by the permutation  $(i(i-1)\dots 1)$  ( $i$  to  $i-1$  etc.)
- ▶ **Example** Above we get  $(4321)$  and  $(21)$

# Cayley graphs of permutations



- ▶ Graph  $\Gamma$  Vertices = permutations, edges = top-to-random shuffles ( $i(i-1)\dots 1$ )
- ▶ Key idea Model top-to-random shuffles as a random walk on  $\Gamma$

## Enter, the theorem

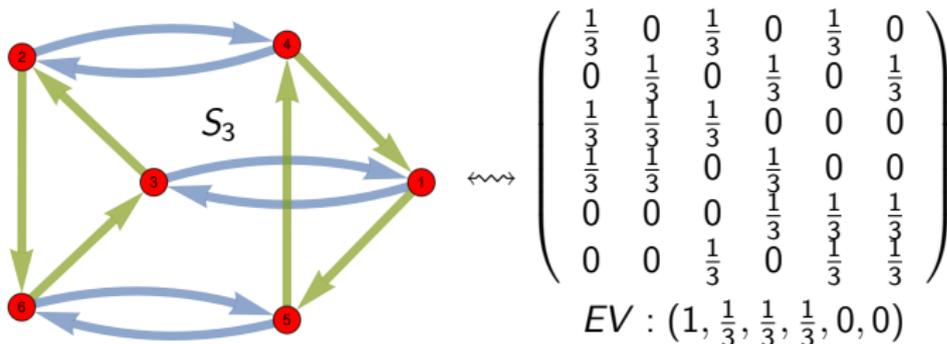
Random walk on  $\Gamma$  (+ loops for “id shuffle”) with all shuffles equally likely has:

- (i) The walk is ergodic **The deck will eventually be mixed**
- (ii) For  $c > 0$  the distance from the uniform distribution after  $k$  shuffles is

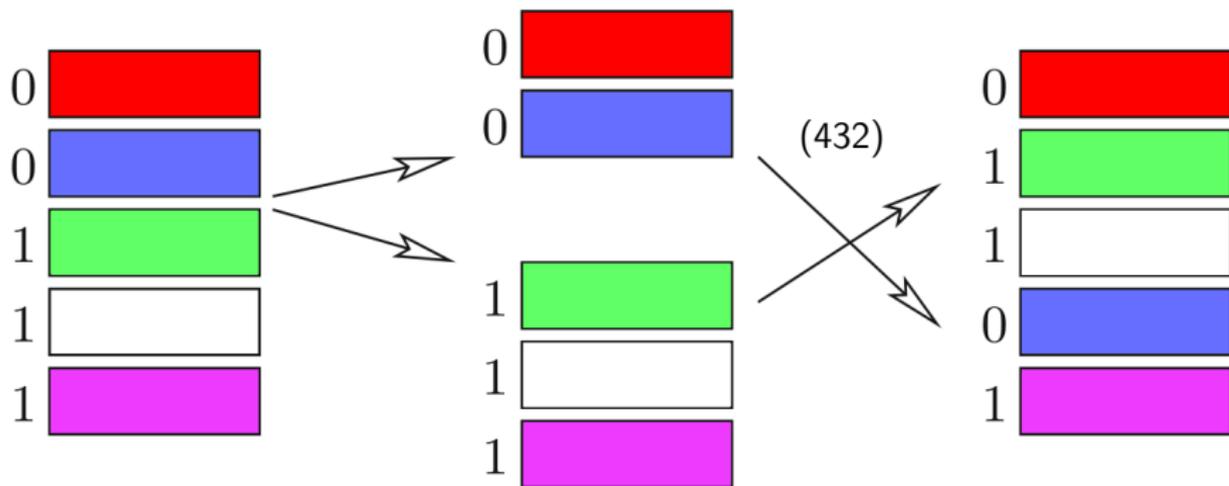
$$|\text{top-to-random}_k - \text{uniform}| \leq e^{-c}$$

where  $k \geq n \log n + cn$  for  $n$  cards

- ▶ For  $n = 52$  and  $c = 5$  one needs  $\approx 465$  shuffles **Inefficient but gets you there**
- ▶ The proof involves analyzing the spectrum of  $\Gamma$  via rep theory



## Riffle shuffle random walk



- ▶ The riffle shuffle can also be modeled by a random walk
- ▶ The edges are now permutation with exactly two rising sequences
- ▶ This walk is also ergodic The deck will eventually be mixed
- ▶ One gets to the mixed state much faster than for the top-to-random shuffle

**Thank you for your attention!**

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I hope that was of some help.